METHODS, SYSTEMS AND COMPUTER PROGRAM PRODUCTS FOR CONTENT-BASED ROUTING VIA ACTIVE TCP CONNECTION TRANSFER

Field of the Invention

The present invention relates to network communications and more particularly to network
5   communications to a cluster of data processing systems.

Background of the Invention

As the use of the Internet has increased, in general, so has the demand placed on servers on the
10   Internet. One technique which has been used to address this increase in demand has been through the use of multiple servers which perform substantially the same function. Applications, such as Telnet or Internet Mail Access Protocol (IMAP)/ Post Office Protocol 3 (POP3)
15   mail serving, may need to connect a Transmission Control Protocol (TCP) client to a particular TCP server of a set of similar, but not identical, servers. However, the particular server instance typically cannot be selected until after information from the client has been
20   received. For example, an Internet Service Provider (ISP) may have multiple e-mail servers to which users may connect to obtain their mail. However, when multiple servers which perform substantially the same function are

present, selecting which server a user should be connected to may present difficulties.

In the e-mail example discussed above, one conventional approach has been to have users individually configure each client application to request a connection to a dedicated server. Such may be accomplished by providing each server with a unique name or Internet Protocol address and configuring the client to specify the name or address when making a connection. Such approaches may, however, present difficulties in maintaining balanced workload between the servers as users come and go. Furthermore, reconfiguring a large population of client applications may present administrative difficulties.

Another approach to routing clients to specific servers has been through an application unique protocol between the client and the server which performs application redirection. In application redirection, a client typically establishes a first connection to a first server which sends a redirect instruction to the client. Upon receiving the redirection instruction, the client disconnects from the initial server and establishes a second connection to the specified server. One difficulty with such an approach, however, is that the client and the server typically must implement the application-unique protocol to provide the redirection and, thus, the redirection is not transparent to the client.

Another approach is known as proxying, where the client establishes an initial connection to a proxy application and the proxy application forms a second connection with the proper server after obtaining enough information from the client to select a server. Such an approach may have the advantage that the selection and communication with the selected server by the proxy may

be transparent to the client. However, both inbound and outbound communications must, typically, traverse a protocol stack twice to be routed by the proxy application. First, the communications traverse the protocol stack to the proxy application and again traverse the protocol stack when routed by the proxy application. Such traversals of the protocol stack may consume significant processing resources at the server executing the proxy application.

In additional approaches, the client establishes a connection to a proxy, which in turn establishes a connection to the ultimate server. Either at a low level in the stack or by instructing an external router, a TCP connection translation function is set up which causes the router or stack to perform modification on all incoming and outgoing TCP segments. The modifications may include the server-side address (destination for incoming requests, source address for outgoing replies) in the IP header, sequence numbers in the TCP header, window sizes, and the like. Such an approach may not require traversal of the entire TCP stack, but may result in every TCP segment requiring modification, and if IP addresses flow in the application data, the connection translation function will not translate such addresses unless specifically programmed to scan all the application data. This approach also requires all flows for the connection, both inbound and outbound, to traverse a single intermediate node, making it a single point of failure (like the proxy).

Furthermore, the Locality-Aware Request Distribution system developed at Rice University is described as providing content-based request distribution which may provide the ability to employ back-end nodes that are specialized for certain types of requests. A "TCP handoff protocol" is described in which incoming requests

are "handed off to a back-end in a manner transparent to a client, after the front-end has inspected the content of the request." *See* Pai *et al.*, "Locality-Aware Request Distribution in Cluster-based Network Servers",
5   Proceedings of the 8th International Conference on Architectural Support for Programming Languages and Operating Systems, San Jose, CA, October, 1998. *See also*, Aron *et al.*, "Efficient Support for P-HTTP in Cluster-Based Web Servers", Proceedings of the USENIX
10   1999 Annual Technical Conference, Monterey, CA, June, 1999.

Other approaches to moving a client connection or session from one server to another include Virtual Telecommunications Access Method (VTAM) multi-node
15   persistent session support. VTAM multi-node persistent session support allows for recovering a System Network Architecture (SNA) session state on another VTAM when an application fails and is restarted. However, typically, a client must re-authenticate to the applications or
20   other system using multi-node persistent sessions. Furthermore, such a movement from a first VTAM to a second VTAM typically only occurs after a failure.

VTAM also supports CLSDEST PASS, which causes one SNA session with the client to be terminated and another
25   initiated without disrupting the application using the sessions. Such a movement from one session to another, however, typically requires client involvement.

Summary of the Invention

30   Methods, systems and computer program products according to embodiments of the present invention provide for distributing Transmission Control Protocol (TCP) connections to a specific data processing system in a cluster of data processing systems by establishing a TCP
35   connection between a client and a first data processing

system in the cluster of data processing systems and obtains information through a plurality of request/response communications with the client over the TCP connection. The information obtained over the TCP connection to the first data processing system is evaluated to select a target data processing system in the cluster of data processing systems for the TCP connection. The TCP connection is transferred from the first data processing system to the selected target data processing system so that the transfer of the TCP connection is transparent to the client.

In further embodiments of the present invention, transferring the TCP connection from the first data processing system to the selected target data processing system may be accomplished by providing connection state information associated with the connection to the selected target data processing system and routing subsequent communications associated with the TCP connection to the selected target data processing system without creating a second TCP connection.

Furthermore, evaluation of the information received over the TCP connection to the first data processing system may be accomplished by providing information received over the TCP connection to an application executing on the first data processing system. The application executing on the first data processing may then evaluate the provided information to select a target instance of an application executing on a target data processing system. The transfer of the TCP connection from the first data processing system to the selected target data processing system may be provided by transferring the TCP connection to the selected target instance of the application executing on the target data processing system.

In still further embodiments of the present invention, the selected target instance of the application executing on the target data processing system may be notified of a TCP connection to be transferred to the selected target instance of the application. A confirmation of acceptance of the transfer of TCP connection may be provided by the selected target instance of the application. The TCP connection is transferred to the selected target instance of the application executing on the target data processing system if the confirmation of acceptance indicates that the selected target instance of the application accepts the transfer of the TCP connection. In particular embodiments of the present invention, the target instance of the application executing on the target data processing system may be a web server.

In still further embodiments of the present invention, the application executing on the first data processing system determines application state information based on the provided information and the selected target instance of the application. The application state information is then provided to the selected target instance of the application executing on the target data processing system. The selected target instance of the application may receive the application state information and establish a state of the target instance of the application based on the received application state information such that the transfer to the target instance of the application is transparent to the client.

Furthermore, the application executing on the first data processing system and the target instance of the application executing on the target data processing system may be instances of the same application.

Alternatively, the application executing on the first data processing system may be a routing application.

Additionally, the connection may be transferred from the selected data processing system to a second selected data processing system. Thus, multiple transfers of the connection may be provided.

In particular embodiments of the present invention, the obtained information is obtained by peeking at information provided over the TCP connection, and the obtained information is provided to the selected data processing system as part of the transfer.

In additional embodiments of the present invention, methods, systems and computer program products are provided for transferring a Transmission Control Protocol (TCP) connection to a specific data processing system in a cluster of data processing systems, the cluster of data processing systems having an associated dynamically routable virtual Internet Protocol address (DVIPA). A connection utilizing the DVIPA is established between a client and a routing application utilizing a routing communication protocol stack at a first data processing system in the cluster of data processing systems. The routing application obtains information from the client over the connection to the routing application and selects a target application for transfer of the connection based on the received information. The routing communication protocol stack is notified of the selected target application and sends a connection transfer message to a target communication protocol stack associated with the selected target application. The connection transfer message contains connection state information associated with the connection to the routing application. Subsequently received communications over the connection are routed to the target communication protocol stack. The target communication protocol stack

notifies the target application of the transfer of the connection to the target application and sets a state of a connection to the target application to the state specified by the connection state information associated with the connection to provide a transferred connection. The target application then communicates with the client utilizing the transferred connection.

In further embodiments of the present invention, the routing application provides application state information to the routing communication protocol stack. The application state information specifies a state of the selected target application based on the communications with the client. The routing communication protocol stack provides the application state information to the target communication protocol stack which provides the application state information to the target application. The target application resumes communications with the client from the application state specified by the provided application state information.

In additional embodiments of the present invention, the target application sends a response message to the routing application. The response message indicates whether the target application accepts or rejects the transfer of the connection. The routing application also closes a socket associated with the connection if the response message from the target application indicates that the target application accepts the transfer of the connection. The routing application may also select a second target application if the response message rejects the transfer of the connection. The routing communication protocol stack is notified of the selection of the second target application so as to initiate transfer of the connection to the second selected target application.

In still further embodiments of the present invention, the routing application sends an error message to the client over the connection if the response message indicates that the transfer of the connection is

5   rejected.

In particular embodiments of the present invention, the routing application opens a control socket to the routing communication protocol stack so as to allow bi-directional communication between the routing

10   communication protocol stack and the routing application. The target application may also open a control socket to the target communication protocol stack to allow bi-directional communication between the target application and the target communication protocol stack. The routing

15   application may be identified to the routing communication protocol stack as a routing application when the routing application opens the control socket. The control sockets may be User Datagram Protocol (UDP) sockets.

20   The routing communication protocol stack may provide the routing application with an identification of potential target applications listening to the DVIPA utilizing the control socket. Also, the routing application may be provided with updated identifications

25   of potential target applications listening to the DVIPA utilizing the control socket.

In still further embodiments of the present invention, the cluster of data processing system includes a plurality of communication protocol stacks and a

30   corresponding plurality of associated applications listening to the DVIPA. In such embodiments, the plurality of communication protocol stacks may notify the routing communication protocol stack if an application is listening to the DVIPA.

In particular embodiments of the present invention, the cluster of data processing systems is a SYSPLEX cluster.

Additionally, transfer of the connection by the routing communication protocol stack may include updating a connection routing table (CRT) associated with the routing communication protocol stack to route communications to the connection to the target communication protocol stack.

As will further be appreciated by those of skill in the art, the present invention may be embodied as methods, apparatus/systems and/or computer program products.

Brief Description of the Drawings

Figure 1 is block diagram of a system suitable for content-based transfer of active TCP connections according to embodiments of the present invention;

Figure 2 is a flowchart illustrating operations for content-based transfer of active TCP connections according to embodiments of the present invention;

Figure 3 is block diagram of a cluster of data processing systems incorporating embodiments of the present invention;

Figure 4 is a flowchart illustrating initialization operations of protocol stacks and applications according to embodiments of the present invention; and

Figure 5 is a flowchart illustrating operations according to further embodiments of the present invention.

Detailed Description of the Invention

The present invention now will be described more fully hereinafter with reference to the accompanying

drawings, in which preferred embodiments of the invention are shown. This invention may, however, be embodied in many different forms and should not be construed as limited to the embodiments set forth herein; rather,

5   these embodiments are provided so that this disclosure will be thorough and complete, and will fully convey the scope of the invention to those skilled in the art. Like numbers refer to like elements throughout.

       As will be appreciated by those of skill in the

10   art, the present invention can take the form of an entirely hardware embodiment, an entirely software (including firmware, resident software, micro-code, *etc*.) embodiment, or an embodiment containing both software and hardware aspects. Furthermore, the present invention can

15   take the form of a computer program product on a computer-usable or computer-readable storage medium having computer-usable or computer-readable program code means embodied in the medium for use by or in connection with an instruction execution system. In the context of

20   this document, a computer-usable or computer-readable medium can be any means that can contain, store, communicate, propagate, or transport the program for use by or in connection with the instruction execution system, apparatus, or device.

25   The computer-usable or computer-readable medium can be, for example, but is not limited to, an electronic, magnetic, optical, electromagnetic, infrared, or semiconductor system, apparatus, device, or propagation medium. More specific examples (a nonexhaustive list) of

30   the computer-readable medium would include the following: an electrical connection having one or more wires, a removable computer diskette, a random access memory (RAM), a read-only memory (ROM), an erasable programmable read-only memory (EPROM or Flash memory), an optical

35   fiber, and a portable compact disc read-only memory (CD-

ROM). Note that the computer-usable or computer-readable medium could even be paper or another suitable medium upon which the program is printed, as the program can be electronically captured, via, for instance, optical

5   scanning of the paper or other medium, then compiled, interpreted, or otherwise processed in a suitable manner if necessary, and then stored in a computer memory.

Embodiments of the present invention provide for content-based transfers of an active TCP connection from

10   one data processing system to another. Such a transfer of an active connection may be beneficial as it may allow for selection of a particular data processing system for handling a client's requested interaction without the client knowing that the transfer has taken place. Such a

15   transfer may include transferring both application state information and connection state information. By transferring the existing active connection and application state, rather than establishing a new connection, the client need not repeat providing the

20   initial information to the data processing system to which the connection is transferred. Furthermore, by transferring the connection to a new data processing system, rather than merely routing packets to and from the new data processing system over a separate

25   connection, administrative overhead may be reduced and scalability may be improved.

A system suitable for use in embodiments of the present invention is illustrated in **Figure 1**. As seen in **Figure 1**, a client **10** communicates with one or more of a

30   plurality of data processing systems **12**. The plurality of data processing systems **12** may include a first data processing system **20**, which functions as a routing data processing system, and several additional data processing systems **22, 22'** and **22''**. The additional data processing

systems **22**, **22'** and **22''** function as server data

processing systems as they are executing instances of an

application from which the client **10** is requesting

interaction. To request interaction with the server data

5   processing systems, the client **10** establishes an initial

TCP connection **30** with the first data processing system

20 of the plurality of data processing systems **12**.

The first data processing system **20** communicates

with the client **10** over the initial connection **30**, for

10   example, through a plurality of request/response

exchanges with the client, until such time as sufficient

information is received that the first data processing

system **20** can determine which of the additional data

processing systems **22**, **22'** and **22''** is to receive the

15   connection with the client **10**. When such a determination

is made, the first data processing system **20** provides

connection and/or application state information to the

selected data processing system.

In the example illustrated in **Figure 1**, the first

20   data processing system **20** has selected the data

processing system **22** and provides the connection and/or

state information to the selected data processing system

22 over the communication path **32**. The communication

path **32** may be any type of communication path suitable

25   for transferring information between the first data

processing system **20** and the selected data processing

system **22** and may include, for example, a TCP connection,

a UDP connection or a connection through a cross-coupling

facility of a SYSPLEX (and XCF connection). If accepted

30   by the selected data processing system **22**, the selected

data processing system **22** continues to communicate with

the client **10** over the transferred connection **30'**, which

is the initial connection **30** transferred to the selected
data processing system **22**.

Operations for transferring an active connection are
illustrated in further detail in **Figure 2**. As seen in
Figure 2, the first data processing system **20** receives a
connection request from the client **10** (block **50**). In
response to the request, the first data processing system
20 establishes the connection **30** to the client **10** (block
52). The first data processing system **20** communicates
with the client **10** over the connection **30** to obtain
information from the client **10** through a plurality of
request/response communications until sufficient
information is obtained to allow the first data
processing system **20** to select one of the other data
processing systems **22**, **22'** and **22''** (block **54**). For
example, in certain embodiments of the present invention,
the information may be a user name and/or password from
which a mail server is selected. Similarly, the user may
be authenticated as part of the selection process.

In any event, once the information is received from
the client **10**, the first data processing system **20**
selects one of the other servers **22**, **22'** and **22''** for
transfer of the initial connection **30** based on the
received information (block **56**). Connection state
information and, optionally, application state
information is collected by the first data processing
system **20** and provided to the selected data processing
system (block **58**), for example, the data processing
system **22** of **Figure 1**. The particular application state
information may vary from application to application.
The connection state information should be sufficient to
allow the initial connection **30** to be transferred to the
selected data processing system **22** to provide the

transferred connection 30' without disruption of the
initial connection 30. For example, for a TCP
connection, the state information may include the source
and destination IP addresses and port numbers, the

5   received and sent packet sequence numbers, the client and
routing application advertised windows, the negotiated
maximum segment size and/or scaling parameters, if any
etc.

       The state of the initial connection 30 may also be
10   "frozen" until the initial connection 30 is transferred
such that any subsequent packets received by the first
data processing system 20 may be discarded by the first
data processing system 20 so as to invoke retransmission
by the client 10. Alternatively, such packets could be
15   buffered and forwarded to the selected data processing
system 22 for acknowledgment by the selected data
processing system 22 over the transferred connection 30'.

       After the application state and/or connection state
information is provided to the selected data processing
20   system 22, the selected data processing system 22 takes
over the initial connection 30 to provide the transferred
connection 30' at the state specified by the connection
state information and the application may continue from
the state specified by the application state information
25   (block 60). The selected data processing system 22 may
communicate directly wit the client 10 over the
transferred connection 30' which corresponds to the
initial connection 30 taken over by the selected data
processing system.

30   As will be appreciated by those of skill in the art
in light of the above discussion, from the perspective of
the client, the initial connection 30 and the transferred
connection 30' are the same connection. Thus, content-
based routing may be achieved and the active connection

transferred to a selected data processing system in a manner which is transparent to the client **10**.

In particular embodiments of the present invention systems, methods, and/or computer program products are provided which allow for a single IP address being associated with a plurality of communication protocol stacks in a cluster of data processing systems by providing a routing protocol stack which associates a Virtual IP Address (VIPA) and port with other communication protocol stacks in the cluster and routes communications to the VIPA and port to the appropriate communication protocol stack. VIPAs capable of being shared by a number of communication protocol stacks are referred to herein as "dynamic routable VIPAs" (DVIPA). While the present invention is described with reference to a specific embodiment in a System/390 Sysplex, as will be appreciated by those of skill in the art, the present invention may be utilized in other systems where clusters of computers utilize virtual addresses by associating an application or application group, rather than a particular communications adapter with the addresses. Thus, the present invention should not be construed as limited to the particular exemplary embodiments described herein.

In particular, embodiments of the present invention may be utilized in systems, such as those described in commonly assigned United States Patent Application Serial No. 09/640,409, entitled "METHODS, SYSTEMS AND COMPUTER PROGRAM PRODUCTS FOR CLUSTER WORKLOAD DISTRIBUTION" (Attorney Docket No. 5577-205), United States Patent Application Serial No. 09/640,438, entitled "METHODS, SYSTEMS AND COMPUTER PROGRAM PRODUCTS FOR FAILURE RECOVERY FOR ROUTED VIRTUAL INTERNET PROTOCOL ADDRESSES" (Attorney Docket No. 5577-206), United States Patent Application Serial No. 09/640,412, entitled "METHODS,

SYSTEMS AND COMPUTER PROGRAM PRODUCTS FOR NON-
DISRUPTIVELY TRANSFERRING A VIRTUAL INTERNET PROTOCOL
ADDRESS BETWEEN COMMUNICATION PROTOCOL STACKS" (Attorney
Docket No. 5577-207), and United States Patent

5   Application Serial No. 09/401,419, entitled "METHODS,
SYSTEMS AND COMPUTER PROGRAM PRODUCTS FOR AUTOMATED
MOVEMENT OF IP ADDRESSES WITHIN A CLUSTER," the
disclosures of which are incorporated herein by reference
as if set forth fully herein.

10   A cluster of data processing systems is illustrated
in **Figure 3** as a cluster of nodes in Sysplex **112**. As
seen in **Figure 3**, the client **10** may initially communicate
with a routing protocol stack **120** which routes
communications to a particular DVIPA. The DVIPA may be

15   specified as a DVIPA for content-based routing by
including in the VIPADISTRIBUTE statement of the DVIPA a
configuration parameter which specifies that the DVIPA is
to be content routed. For example, the configuration
parameter may be "CONTENTROUTED" which would specify that

20   connection requests to the DVIPA of the VIPADISTRIBUTE
statement are to be initially routed to the routing
application **124** rather than being distributed to the
other data processing systems as would be the case, for
example, for a conventional connection request to a

25   DVIPA.

        As is further illustrated in **Figure 3**, the routing
communication protocol stack **120** has associated with it a
control socket **122** over which bi-directional
communications between the routing communication protocol

30   stack **120** and the routing application **124** are carried
out. The control socket **122** may, for example, be a User
Datagram Protocol (UDP) socket opened by the routing
application **124**. When the UDP socket is opened, the
routing application **124** may identify itself to the

routing communication protocol stack **120** such that the
routing communication protocol stack **120** may associate
the routing application **124** with the DVIPA.

   The routing communication protocol stack **120**
communicates with potential target communication protocol
stacks **130**, **130'** for content-based routing in the cluster
112 utilizing the cross-coupling facility (XCF) **126** and
XCF messages.   In embodiments of the present invention,
the routing application **124** may also communicate with
target applications, such as applications **134** and **134'**,
utilizing messages to its control socket **122** which
results in XCF messages being transmitted to the
potential target communication protocol stacks **130** which
may provide corresponding messages to the applications
134 and **134'** through the control sockets **132**.

   The potential target applications **134** and **134'** may
identify themselves to the communication protocol stacks
130 when the control sockets **132** are opened.   The
identification of the potential target applications **134**
and **134'** may be provided to the routing communication
protocol stack **120** as part of the DVIPA messages which
are distributed throughout the cluster **112**.   Such an
identification may, for example, identify the
applications **134** and **134'** by, for example, the operating
system name and/or jobname for the applications **134** and
134'.   Furthermore, communications from the routing
application **124** to the communication protocol stacks **130**
and/or the potential target applications **134** and **134'** may
be sent utilizing the dynamic XCF address of the
communication protocol stacks **130**.

   Also illustrated in **Figure 3** is a conventional
application **136** which does not utilize a control socket
and communicates utilizing the communication protocol

stack **130'**. Because the communication protocol stack
130' does not have a control socket, the communication
protocol stack **130'** may be a potential target for a
conventional DVIPA or may be a potential target for a
content-based routing DVIPA which does not require
application state information to be passed to the
application **136**.

Operations for the initialization of the
communication protocol stacks and applications according
to embodiments of the present invention are illustrated
in **Figure 4**. As seen in **Figure 4**, the target
communication protocol stacks **130** and **130'** are started
and the routing communication protocol stack **120** is also
started (block **200**). The potential target applications
134, **134'** and **136** are also started and, if the
application is to receive application state information,
a UDP socket is opened on the communication protocol
stack **130** and the application identified to the
communication protocol stack **130** (block **202**). The target
communication protocol stack **130** may, however, block
accepts to the listening socket of the potential target
applications **134** and **134'** until a connection is
transferred to the applications.

The routing communication protocol stack **120** is
notified of the applications which are listening on the
DVIPA which is defined as supporting content-based
routing (block **204**). Such notification may be carried
out by distributing the information between the
communication protocol stacks **120**, **130** and **130'**, for
example, using the XCF **126**, and may include an
identification of the particular application listening to
the DVIPA as is described above. Content-based routing,
however, will not be provided until a routing application

is started and identified to the routing protocol stack **120**.

The routing application **124** is also started on the routing communication protocol stack **120** (block **206**). The routing application **124** establishes a UDP socket as its control socket **122** and identifies itself to the routing communication protocol stack **120** as a routing application (block **208**), for example, by issuing a getibmsockopt(), and identifying the DVIPA and port, as well as the socket to which subsequent client requests should be routed. The routing application **124** may identify itself as associated with a specific DVIPA and port or may identify its listening socket to the routing communication protocol stack **120**.

Because the routing application **124** is identified as a routing application, the routing communication protocol stack **120** also provides the target application identification to the routing application **124** over the control socket **122** (block **210**). Such information may be automatically provided to the routing application **124** or may be provided in response to a request from the routing application **124**. However, once provided, the information may be updated as application instances are terminated or are started and bind to the communication protocol stacks in the cluster **112**. Through the use of such an update, the routing communication protocol stack **120** may provide current information to the routing application **124** as the information is provided to the routing communication protocol stack **120**. Such information may include Destination Port Table information which identifies communication protocol stacks having applications listening to the IP address and port of the DVIPA, as well as job names of applications listening to the IP address and port.

Operations for transferring an active connection
will now be further described with reference to **Figure 5**.
As seen in the embodiments illustrated in **Figure 5**, the
routing communication protocol stack **120** receives a

5   connection request at the listening socket of the routing
application **124** (or applications using port sharing) or
the IP address and port of the DVIPA for which content-
based routing is specified (block **220**). Because the
connection request is to the IP address and the port of

10   the DVIPA, the routing communication protocol stack **120**
establishes a connection to the routing application **124**
for the connection request (block **222**).

The routing application evaluates the data from the
connection request (block **224**) and determines if

15   sufficient information has been provided by the client to
make a routing determination (block **226**). If sufficient
information has not been provided to make a routing
determination (block **226**), the routing application **124**
continues communications with the client **10** (block **228**),

20   evaluates the provided information (block **224**) and again
determines if sufficient information has been received to
make the routing determination (block **226**). These
operations continue until enough information is received
from the client **10** to make the routing determination. In

25   certain embodiments, the information provided by the
client may be "peeked" at by the routing application such
that the information may be evaluated without the routing
application **124** "receiving" the information.

When the routing application **124** has obtained

30   sufficient information from the client **10** to make the
routing determination (block **226**), the routing
application **124** notifies the routing communication
protocol stack **120** to request transfer of the connection

to the target application (block **230**). The routing
communication protocol stack **120** builds a transfer
request by, for example, incorporating the connection
state information into a connection transfer message
5   (block **232**) and the transfer request is sent as an XCF
message to the target stack (block **234**).

The routing application **124** may request the transfer
of the active connection to the selected target
application **134** by sending the request to the routing
10   communication protocol stack **120** over the control socket
122 (block **230**). The transfer request message to the
routing communication protocol stack **120** is preferably of
similar format to an XCF message to the selected target
communication protocol stack such that only minimal
15   modification of the message need be carried out by the
routing communication protocol stack **120** to send the
message to the target application **134**. Preferably, the
routing application waits to request the transfer until a
time when communications with the client **10** have quiesced
20   or until the client **10** is waiting for a response. The
request to transfer may include an identification of the
selected target application and the application state
information, if any. Furthermore, if the routing
application only peeks at the information from the
25   client, the information from the client may also be
provided. Alternatively, the information could be
provided after the connection is transferred.

The routing communication protocol stack **120** may
reject the request if, for example, the request specifies
30   an unknown destination for the transfer or if data is
queued for transmission to the client **10**. If the request
to transfer the connection is accepted, the connection
state may be frozen so that the connection may be resumed
by the selected target application and target

communication protocol stack. Such a freeze of the
connection state may be accomplished, for example, by the
routing communication protocol stack **120** ceasing sending
acknowledgments to datagrams from the client **10**. Any
5   received datagrams may be disregarded or may be buffered
and provided to the selected target application through
the selected target communication protocol stack.

As described above, the application state may depend
on the particular application being utilized. For
10   example, the application state may indicate that a user
has been authenticated and that the application may
resume operations after the authentication. In such
embodiments, the application may be viewed as distributed
across the cluster **112** such that initial, common,
15   portions of the application are carried out by the
routing application **124** which then transfers operations
to the selected target application, such as application
134. Alternatively, for "stateless" applications, *i.e.*
applications which always assume a predefined state for
20   all operations, such as a conventional web server which
serves web pages in response to Hypertext Transport
Protocol (HTTP) requests, the routing application **124** may
merely transfer control of the connection to the target
communication protocol stack and the target application.
25   In such embodiments, application state information need
not be provided to the target application.

Returning to **Figure 5**, the routing communication
protocol stack sends the transfer request message to the
communication protocol stack associated with the selected
30   target application **134** (block **234**). As described above,
the transfer request message may be sent as an XCF
message which specifies the selected target application
as the destination, by including, for example, the
dynamic XCF address of the communication protocol stack

associated with the selected target application **134** and the jobname of the selected target application **134** in the XCF message.

The target communication protocol stack **130** receives the transfer request message, removes the connection state information from the message and provides the application state information and information sufficient to allow the selected target application **134** to associate the connection with the application state information, such as, for example, client IP address and port, to the selected target application **134** over the control socket **132** (block **236**). The request may also be rejected by the target stack if, for example, the specified target application does not exist.

The target application **134** evaluates the request, for example, by evaluating the application state information and/or other information in the request, to determine if the target application **134** will accept the request (block **238**). If the request is not accepted (block **238**), the target application **134** provides a "rejected" response to the communication protocol stack **130** associated with the target application **134** over the control socket **132** (block **240**). If the request is accepted (block **238**), the target application **134** takes ownership of the connection and sets its application state to the application state specified in the received transfer request, if any. The target application **134** also provides an "accepted" response to the communication protocol stack **130** associated with the target application **134** over the control socket **132** (block **242**).

The target communication protocol stack **130** receives the response from the target application **134** and, if the request is accepted, unblocks the listening port of the selected target application **134** which results in a socket

being established for the connection, sets the connection
state associated with the socket to the received
connection state and reports the response to the routing
communication protocol stack **120** (block **244**). If the
request is not accepted, the target communication
protocol stack may discard the connection state
information and report the response to the routing
communication protocol stack **120**. The routing
communication protocol stack **120** receives the response to
the transfer request and provides the response to the
routing application **124** (block **246**). If the response
indicates that the request was accepted, the routing
communication protocol stack **120** may begin routing
incoming messages to the target communication protocol
stack **130** in the same manner as other DVIPAs (*i.e.*
setting an entry in the CRT so that messages to the DVIPA
from the source IP address are forwarded to the selected
target stack). The routing application **124** may close the
socket associated with the connection when the acceptance
of the transfer is received (block **246**).

　　If the request is rejected, the routing application
124 may select a new target application or send a
rejection to the client **10**. As such, if the request is
rejected, the routing communication protocol stack **120**
may unfreeze the connection to the routing application
124 and allow further communications with the client **10**.
Such communication may include, for example, further
communications to select a target application or error
notifications. If a new target application is selected,
then operations may be repeated beginning with block **232**.

　　While the operations of **Figure 5** illustrate
embodiments of the present invention which provide a
"handshake" between the routing application and the
target application, in alternative embodiments of the

present invention, such a handshake may be omitted. In
such embodiments, the operations of blocks **232, 234, 236,
238, 240, 242, 244** and **246** may be eliminated or modified
such that the routing application sends the application

5   state information, if any, and the connection state
information to the target application and the target
communication protocol stack which unconditionally accept
the transfer of the active connection. If the routing
application may select the target application based on

10   the initial TCP data from the client, then the
application state information need not be provided to the
target application and the unconditional transfer may
occur without the use of a control socket to provide
application state information to the target application.

15   Additionally, if the full handshake is not utilized,
the givesocket() and takesocket() functions, which
conventionally only allow for transfers of ownership of
sockets between applications on a single TCP stack, could
be modified to allow for transfers across TCP stacks in a

20   cluster. The givesocket()/takesocket() may be modified
to provide a giveconnection() and takeconnection()
function where, in addition to the client identification,
for giveconnection() and takeconnection() reserved fields
may be modified to specify the target stack and MVS image

25   by name. Giveconnection() can be expanded to include an
optional "application state" data field which may include
application state data or a pointer to an application
state data block. The contents of the application state
field may be passed to the target application on its

30   control socket, along with the client identification of
the "giving" application. Logical File Support (LFS) for
takeconnection() may be required to create a new vnode
and socket on the receiving stack, and an internal Input
Output Control (IOCTL) function may be defined for the

35   LFS to notify the stack of the taking application socket.

Furthermore, while embodiments of the present
invention have been described with reference to transfer
of active connections between a first system and a second
system, as will be appreciated by those of skill in the
5   art in light of the present disclosure, the active
connection could be further transferred from the second
system to a third system. Thus, for example, particular
functions or operations of an application could be
distributed across multiple data processing systems where
10   the data processing systems transferred the active
connection to the next data processing system upon
completion of the function or operation assigned to the
data processing system. Thus, the present invention
should not be construed as limited to a single transfer
15   of an active connection but may include multiple
transfers.

With regard to error recovery, error recovery may be
provided by the recovery mechanism described in the
above-referenced United States Patent Applications.
20   However, if a routing application is started on a backup
protocol stack, the listening port of the routing
application should be blocked until the backup protocol
stack becomes a primary protocol stack.

Embodiments of the present invention have been
25   described with reference to **Figures 1** through **5** which are
block diagrams and/or flowchart illustrations. It will be
understood that each block of the flowchart illustrations
and/or block diagrams, and combinations of blocks in the
flowchart illustrations and/or block diagrams, can be
30   implemented by computer program instructions. These
program instructions may be provided to a processor to
produce a machine, such that the instructions which
execute on the processor create means for implementing
the functions specified in the flowchart and/or block
35   diagram block or blocks. The computer program

instructions may be executed by a processor to cause a
series of operational steps to be performed by the
processor to produce a computer implemented process such
that the instructions which execute on the processor
5   provide steps for implementing the functions specified in
the flowchart and/or block diagram block or blocks.

Accordingly, blocks of the flowchart illustrations
and/or block diagrams support combinations of means for
performing the specified functions, combinations of steps
10   for performing the specified functions and program
instruction means for performing the specified functions.
It will also be understood that each block of the
flowchart illustrations and/or block diagrams, and
combinations of blocks in the flowchart illustrations
15   and/or block diagrams, can be implemented by special
purpose hardware-based systems which perform the
specified functions or steps, or combinations of special
purpose hardware and computer instructions.

While the present invention has been described with
20   respect to the content-based routing function as a part
of the communication protocol stack, as will be
appreciated by those of skill in the art, such functions
may be provided as separate functions, objects or
applications which may cooperate with the communication
25   protocol stacks. Furthermore, the present invention has
been described with reference to particular sequences of
operations. However, as will be appreciated by those of
skill in the art, other sequences may be utilized while
still benefitting from the teachings of the present
30   invention. Thus, while the present invention is
described with respect to a particular division of
functions or sequences of events, such divisions or
sequences are merely illustrative of particular
embodiments of the present invention and the present

invention should not be construed as limited to such embodiments.

Furthermore, while the present invention has been described with reference to particular embodiments of the present invention in a System/390 environment, as will be appreciated by those of skill in the art, the present invention may be embodied in other environments and should not be construed as limited to System/390 but may be incorporated into other systems, such as a Unix or other environments, by associating applications or groups of applications with an address rather than a communications adapter. Thus, the present invention may be suitable for use in any collection of data processing systems which allow sufficient communication to all of the systems for the use of dynamic virtual addressing. Accordingly, specific references to System/390 systems or facilities, such as the "coupling facility," "ESCON," "Sysplex" or the like should not be construed as limiting the present invention.

In the drawings and specification, there have been disclosed typical preferred embodiments of the invention and, although specific terms are employed, they are used in a generic and descriptive sense only and not for purposes of limitation, the scope of the invention being set forth in the following claims.